

# EUCLIDEAN ISOSCELES SETS AND TWO-DISTANCE EXTREMAL FUNCTIONS

ABSTRACT. A finite set in Euclidean space is called isosceles if every three of its points span an isosceles triangle. Let  $f(d)$  be the largest size of an isosceles subset of  $\mathbb{R}^d$ , let  $g(d)$  be the largest size of a Euclidean two-distance set in  $\mathbb{R}^d$ , let  $s(d)$  be the largest size of a spherical two-distance set in  $\mathbb{R}^d$ , and let  $h(d)$  be the largest size of an isosceles subset of  $\mathbb{R}^d$  which is not a two-distance set. We prove that, for every  $d \geq 2$ ,

$$h(d) \leq \max\{s(d) + 1, s(d - 1) + 3\},$$

and consequently

$$f(d) = \max\{g(d), s(d) + 1, s(d - 1) + 3\}.$$

The proof combines Ionin's complete decomposition theorem for Euclidean isosceles sets, the classical absolute bounds for Euclidean and spherical two-distance sets, and a polynomial argument for the only possible two-point block case.

## 1. INTRODUCTION

A finite set  $S \subset \mathbb{R}^d$  is called *isosceles* if every three distinct points of  $S$  determine an isosceles triangle. Equivalently, every three-point subset of  $S$  determines at most two nonzero distances. Let

$$f(d) := \max\{|S| : S \subset \mathbb{R}^d \text{ is isosceles}\}.$$

This is the Euclidean isosceles-set problem appearing as Problem 503 in the Erdős Problems collection [3]. The present note determines  $f(d)$  in terms of the standard extremal functions for Euclidean and spherical two-distance sets.

Let

$$g(d) := \max\{|X| : X \subset \mathbb{R}^d \text{ is a Euclidean two-distance set}\},$$

and let

$$s(d) := \max\{|X| : X \subset \mathbb{R}^d \text{ is a spherical two-distance set}\}.$$

Here and throughout, “two-distance” means “at most two nonzero distances”. For  $d \geq 2$  this convention gives the same values of  $g(d)$  and  $s(d)$  as the convention “exactly two distances”, since the known two-distance examples are larger than a simplex.

We also put

$$h(d) := \max\{|S| : S \subset \mathbb{R}^d \text{ is isosceles and is not a two-distance set}\}.$$

Then

$$f(d) = \max\{g(d), h(d)\}.$$

Our main theorem is the following.

**Theorem 1.1.** *For every integer  $d \geq 2$ ,*

$$h(d) \leq \max\{s(d) + 1, s(d - 1) + 3\}.$$

*Consequently,*

$$f(d) = \max\{g(d), s(d) + 1, s(d - 1) + 3\}.$$

*Also  $f(1) = 3$ .*

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The lower bound comes from two simple constructions. If  $X \subset \mathbb{R}^d$  is a spherical two-distance set and  $c$  is the center of a sphere containing  $X$ , then  $X \cup \{c\}$  is isosceles, giving  $f(d) \geq s(d) + 1$ . If  $X$  is a spherical two-distance set in a hyperplane  $H \cong \mathbb{R}^{d-1}$ , and if one adds the center of the sphere together with two symmetric points on the perpendicular line through the center, then one obtains an isosceles set of size  $s(d-1) + 3$ . This second construction always has more than two distances.

The upper bound uses Ionin's complete decomposition theorem for Euclidean isosceles sets. The decomposition reduces the problem to two-distance blocks. The only delicate case is when a non-last block consists of two points. In that case all remaining points lie in the perpendicular bisector of the pair. The main new ingredient is a two-shell polynomial lemma, proved in Section 5, which gives exactly the one-point improvement over Blokhuis's general two-distance bound needed for this pair-block case.

## 2. PRELIMINARIES

For a finite set  $X \subset \mathbb{R}^d$ , write

$$\dim_{\text{aff}} X := \dim(\text{aff } X)$$

for its affine dimension. We use the following notation:

$$U(m) := \frac{m(m+3)}{2}, \quad L(m) := \binom{m+1}{2}, \quad G(m) := \binom{m+2}{2} = U(m) + 1.$$

We set  $G(0) = 1$  and  $s(1) = 2$ . The latter convention is only used in the term  $s(d-1) + 3$  when  $d = 2$ .

The classical Delsarte–Goethals–Seidel and Blokhuis bounds give [2, 1]

$$(1) \quad s(m) \leq U(m), \quad g(m) \leq G(m) \quad (m \geq 1),$$

where the assertions for  $m = 1$  are immediate. The regular-simplex midpoint construction gives

$$(2) \quad s(m) \geq L(m) \quad (m \geq 1).$$

For  $m = 1$  this reads  $2 \geq 1$ . We shall also use without further comment that  $g(m)$  and  $s(m)$  are nondecreasing in  $m$ , which follows by embedding lower-dimensional configurations in a coordinate hyperplane.

We shall use Ionin's complete decomposition theorem in the following form.

**Definition 2.1.** Let  $S$  be a finite metric space. A partition

$$S = S_1 \sqcup \cdots \sqcup S_k$$

is called a *complete decomposition* if

- (i) each  $S_i$  is a two-distance set;
- (ii)  $|S_i| \geq 2$  for  $1 \leq i < k$ , and  $|S_k| \geq 1$ ;
- (iii) if  $1 \leq i < j \leq k$ , then every point of  $S_j$  is the center of a sphere containing  $S_i$ .

**Theorem 2.2** (Ionin). *Every finite isosceles set in Euclidean space admits a complete decomposition. Moreover, if*

$$S = S_1 \sqcup \cdots \sqcup S_k$$

*is a complete decomposition of an isosceles set  $S \subset \mathbb{R}^d$ , then*

$$\dim_{\text{aff}} S \geq \sum_{i=1}^k \dim_{\text{aff}} S_i.$$

*In particular, every non-last block  $S_i$ ,  $i < k$ , is spherical.*

*Reference.* This is Ionin's complete decomposition theorem for Euclidean isosceles sets; see [4, Definition 4.1 and Proposition 4.2]. The last assertion follows from the defining property that every later block consists of centers of spheres containing  $S_i$ .  $\square$

The following elementary inequality is used to count several non-last blocks at once.

**Lemma 2.3.** *Let  $a_1, \dots, a_q$  be positive integers and put  $N = a_1 + \dots + a_q$ . If  $q \geq 2$ , then*

$$\sum_{i=1}^q U(a_i) \leq L(N) + 1.$$

Moreover,

$$\sum_{i=1}^q U(a_i) \leq L(N)$$

unless  $q = 2$  and  $\{a_1, a_2\} = \{1, N - 1\}$ .

*Proof.* A direct calculation gives

$$\begin{aligned} L(N) - \sum_{i=1}^q U(a_i) &= \frac{N(N+1) - \sum_i (a_i^2 + 3a_i)}{2} \\ &= \sum_{1 \leq i < j \leq q} a_i a_j - N. \end{aligned}$$

If  $q = 2$ , this equals

$$a_1 a_2 - a_1 - a_2 = (a_1 - 1)(a_2 - 1) - 1,$$

so it is at least  $-1$ , and it is negative exactly when one of  $a_1, a_2$  is 1.

Assume  $q \geq 3$ . For fixed  $N$ , the sum  $\sum_{i < j} a_i a_j$  is minimized when one part is as large as possible and all other parts are 1. Hence

$$\sum_{i < j} a_i a_j \geq (q-1)(N-q+1) + \binom{q-1}{2}.$$

The right side is at least  $N$  for every  $q \geq 3$ . Therefore  $L(N) - \sum_i U(a_i) \geq 0$  when  $q \geq 3$ , and the lemma follows.  $\square$

### 3. LOWER CONSTRUCTIONS

**Proposition 3.1.** *For every integer  $d \geq 2$ ,*

$$f(d) \geq \max\{g(d), s(d) + 1, s(d-1) + 3\}.$$

Moreover,

$$h(d) \geq s(d-1) + 3.$$

*Proof.* The term  $g(d)$  is immediate, since every two-distance set is isosceles.

Let  $X \subset \mathbb{R}^d$  be a spherical two-distance set of size  $s(d)$ , and let  $c$  be the center of a sphere containing  $X$ . Then  $X \cup \{c\}$  is isosceles: triples contained in  $X$  are already isosceles, and every triple of the form  $\{x, y, c\}$  is isosceles because  $|x - c| = |y - c|$ . Hence  $f(d) \geq s(d) + 1$ .

For the third term, let  $H \cong \mathbb{R}^{d-1}$  be a hyperplane, let  $X \subset H$  be a spherical two-distance set of size  $s(d-1)$ , and let  $c \in H$  be the center of a sphere containing  $X$ . Write  $r = |x - c|$  for  $x \in X$ . Let  $p, q$  be the two points on the line through  $c$  perpendicular to  $H$  such that  $c$  is the midpoint of  $pq$  and

$$|p - c| = |q - c| = r.$$

Set

$$Y := X \cup \{p, c, q\}.$$

Then  $Y$  is isosceles. Indeed, the only nontrivial triples are checked from the identities

$$|x - c| = |p - c| = |q - c| = r, \quad |p - x| = |q - x| = \sqrt{2}r \quad (x \in X).$$

The set  $Y$  has more than two distances, since it realizes  $r$ ,  $\sqrt{2}r$ , and  $2r$ . Therefore  $h(d) \geq s(d-1) + 3$ , and the proposition follows.  $\square$

**Remark 3.2.** The construction  $X \cup \{c\}$  gives the lower bound  $s(d) + 1$  for  $f(d)$ , but it does not necessarily give a lower bound for  $h(d)$ . It may happen that the radius of the sphere containing  $X$  is already one of the two distances determined by  $X$ , in which case  $X \cup \{c\}$  is still a two-distance set.

#### 4. REDUCTION TO A TWO-POINT BLOCK

Let

$$M(d) := \max\{s(d) + 1, s(d-1) + 3\}.$$

The purpose of this section is to show that every large non-two-distance isosceles set has a complete decomposition with a two-point first block and one last block of codimension one.

**Proposition 4.1.** *Let  $d \geq 2$ , and let  $S \subset \mathbb{R}^d$  be an isosceles set which is not a two-distance set. Then either*

$$|S| \leq M(d),$$

or  $S$  has a complete decomposition

$$S = S_1 \sqcup S_2$$

such that  $S_1$  consists of two points and  $\dim_{\text{aff}} S_2 = d - 1$ .

*Proof.* Choose a complete decomposition

$$S = S_1 \sqcup \cdots \sqcup S_k.$$

Since  $S$  is not a two-distance set,  $k \geq 2$ . Put  $q = k - 1$ . For  $1 \leq i \leq q$ , set

$$a_i := \dim_{\text{aff}} S_i.$$

The non-last blocks are spherical and have at least two points, so  $a_i \geq 1$  and

$$|S_i| \leq U(a_i) \quad (1 \leq i \leq q).$$

In particular, a one-dimensional non-last block has exactly two points, since a line meets a sphere in at most two points. Let

$$N := a_1 + \cdots + a_q, \quad e := \dim_{\text{aff}} S_k.$$

By Theorem 2.2,  $N + e \leq d$ .

First suppose  $q = 1$ , and write  $a = a_1$ . If  $a \geq 2$ , then the estimates are standard. If  $e = 0$ , then

$$|S| \leq s(a) + 1 \leq s(d) + 1.$$

If  $e = 1$ , then  $S_k$  has at most three points and  $a \leq d - 1$ , so

$$|S| \leq s(a) + 3 \leq s(d-1) + 3.$$

If  $e \geq 2$ , then

$$|S| \leq U(a) + G(e),$$

and

$$L(a+e) + 1 - U(a) - G(e) = ae - a - e \geq 0,$$

because  $a, e \geq 2$ . Hence

$$|S| \leq L(d) + 1 \leq s(d) + 1.$$

Thus the case  $q = 1$ ,  $a \geq 2$ , is bounded by  $M(d)$ .

It remains in the case  $q = 1$  to consider  $a = 1$ . Then  $S_1$  consists of exactly two points. If  $e \leq d - 2$ , then

$$|S| \leq 2 + G(e) \leq 2 + G(d-2) \leq L(d) + 1 \leq s(d) + 1.$$

Therefore the only unbounded alternative with  $q = 1$  is precisely  $k = 2$ ,  $|S_1| = 2$ , and  $\dim_{\text{aff}} S_2 = d - 1$ .

Now suppose  $q \geq 2$ . If the exceptional alternative in Lemma 2.3 does not occur, then

$$\sum_{i=1}^q |S_i| \leq L(N).$$

If  $e = 0$ , then

$$|S| \leq L(N) + 1 \leq L(d) + 1 \leq s(d) + 1.$$

If  $e = 1$ , then

$$|S| \leq L(N) + 3 \leq L(d-1) + 3 \leq s(d-1) + 3.$$

If  $e \geq 2$ , then

$$|S| \leq L(N) + G(e),$$

and

$$L(N+e) + 1 - L(N) - G(e) = e(N-1) \geq 0.$$

Thus  $|S| \leq s(d) + 1$ .

It remains to handle the exceptional alternative in Lemma 2.3. Then  $q = 2$  and the two non-last dimensions are 1 and  $N-1$ . The one-dimensional block contributes two points. If  $e = 0$ , then

$$|S| \leq 2 + s(N-1) + 1 \leq s(d-1) + 3.$$

If  $e = 1$ , then

$$|S| \leq 2 + U(N-1) + 3.$$

Since  $N \leq d-1$ , one has

$$2 + U(N-1) + 3 \leq L(d) + 1 \leq s(d) + 1.$$

Finally, if  $e \geq 2$ , then

$$|S| \leq 2 + U(N-1) + G(e),$$

and

$$L(N+e) + 1 - (2 + U(N-1) + G(e)) = e(N-1) - 1 \geq 0.$$

Hence  $|S| \leq s(d) + 1$  also in this case. The proposition follows.  $\square$

## 5. THE TWO-SHELL POLYNOMIAL LEMMA

The following lemma is the main point needed for the two-point block case.

**Lemma 5.1.** *Let  $m \geq 1$ , let  $X \subset \mathbb{R}^m$  be a two-distance set with squared distances  $0 < \alpha < \beta$ , and let  $h > 0$ . Suppose that*

$$|x|^2 + h^2 \in \{\alpha, \beta\} \quad (x \in X).$$

Put

$$X_1 := \{x \in X : |x|^2 + h^2 = \alpha\}, \quad X_2 := \{x \in X : |x|^2 + h^2 = \beta\}.$$

If  $|X_1| \geq 2$  and  $|X_2| \geq 2$ , then

$$|X| \leq \binom{m+2}{2} - 1 = U(m).$$

*Proof.* Let  $\eta := h^2$ ,  $r_1 := \alpha - \eta$ ,  $r_2 := \beta - \eta$ , and  $\Delta := \beta - \alpha$ . Since  $|X_1| \geq 2$ , the first shell has positive radius; hence

$$0 < \eta < \alpha.$$

If  $\dim_{\text{aff}} X < m$ , then Blokhuis's bound in the affine span of  $X$  gives

$$|X| \leq G(m-1) \leq G(m) - 1 = U(m),$$

so we may assume that  $X$  has affine dimension  $m$ .

By Blokhuis's bound,  $|X| \leq G(m)$ . We shall exclude equality. Suppose, for contradiction, that

$$|X| = G(m) = \binom{m+2}{2}.$$

Let  $u = (u_1, \dots, u_m)$ , let  $z$  be an additional variable, and set

$$\sigma := u_1^2 + \dots + u_m^2 + z^2, \quad \rho := u_1^2 + \dots + u_m^2.$$

For  $x \in X$ , write  $r_x := |x|^2$ , and define

$$F_x(u, z) := \frac{(\sigma - 2x \cdot u + r_x - \alpha)(\sigma - 2x \cdot u + r_x - \beta)}{\alpha\beta}.$$

Then

$$F_x(y, 0) = \delta_{xy} \quad (x, y \in X).$$

Also set

$$Q(u, z) := (\sigma - r_1)(\sigma - r_2).$$

Since  $|y|^2 \in \{r_1, r_2\}$  for every  $y \in X$ , one has

$$Q(y, 0) = 0 \quad (y \in X).$$

Let  $R \in \mathbb{R}$  be a parameter, and fix a nonzero real number  $\varepsilon$ . Define

$$\Phi_x := F_x + \left( -\frac{1}{\alpha\beta} + \varepsilon(2r_x - R) \right) Q \quad (x \in X).$$

Then

$$\Phi_x(y, 0) = \delta_{xy} \quad (x, y \in X).$$

Each  $\Phi_x$  lies in the vector space

$$W := \text{span}\{\sigma^2, \sigma u_1, \dots, \sigma u_m, u_i u_j \ (1 \leq i \leq j \leq m), \sigma, u_1, \dots, u_m, 1\}.$$

The displayed spanning polynomials are linearly independent. Indeed, terms of different total degree cannot cancel; in degree two, the polynomial  $\sigma$  is independent of the polynomials  $u_i u_j$  because it contains the term  $z^2$ , while no polynomial  $u_i u_j$  does. Hence

$$\dim W = 1 + m + \binom{m+1}{2} + 1 + m + 1 = \binom{m+3}{2}.$$

Since

$$|X| + m + 3 = G(m) + m + 3 = \binom{m+3}{2} + 1,$$

the family

$$\{\Phi_x : x \in X\} \cup \{u_1, \dots, u_m, \sigma, \rho, 1\}$$

is linearly dependent. Thus there are real coefficients  $c_x, a_i, b, c,$  and  $d$ , not all zero, such that

$$(3) \quad \sum_{x \in X} c_x \Phi_x + \sum_{i=1}^m a_i u_i + b\sigma + c\rho + d = 0$$

as a polynomial identity.

Put

$$C := \sum_{x \in X} c_x, \quad V := \sum_{x \in X} c_x r_x x, \quad M := \sum_{x \in X} c_x x x^T.$$

We compare coefficients in (3). The coefficient of  $\sigma^2$  gives, after division by  $\varepsilon$ ,

$$(4) \quad \sum_{x \in X} c_x (2r_x - R) = 0,$$

and hence

$$(5) \quad \sum_{x \in X} c_x r_x = \frac{R}{2} C.$$

The coefficients of the terms  $\sigma u_i$  give

$$(6) \quad \sum_{x \in X} c_x x = 0.$$

Comparing quadratic terms gives

$$(7) \quad M = \mu I_m, \quad \mu = \frac{RC}{2m},$$

where the value of  $\mu$  follows by taking traces and using (5). The linear coefficients give

$$(8) \quad a := (a_1, \dots, a_m) = \frac{4}{\alpha\beta} V.$$

If  $P := b + c$ , then the coefficients of  $z^2$  and  $u_i^2$  give

$$(9) \quad P = \frac{C}{\alpha\beta} \left( 2\eta - \frac{m+2}{m} R \right).$$

Finally, the constant coefficient gives

$$(10) \quad d = \frac{C}{\alpha\beta} (\eta R + \alpha\beta - 2(\alpha + \beta)\eta + 2\eta^2).$$

For completeness, we note that the terms involving  $\varepsilon$  vanish in these comparisons because of (4).

Evaluating (3) at  $(y, 0)$ , with  $y \in X$ , gives

$$(11) \quad c_y + a \cdot y + P|y|^2 + d = 0.$$

Let

$$C_2 := \sum_{y \in X_2} c_y, \quad V_2 := \sum_{y \in X_2} c_y y.$$

From (5) and (6),

$$(12) \quad C_2 = \frac{(R/2 - r_1)C}{\Delta}, \quad V_2 = \frac{V}{\Delta}.$$

Now multiply (11) by  $c_y$  and sum over  $y \in X_2$ . Using (8) and (12), we obtain

$$(13) \quad 0 = \sum_{y \in X_2} c_y^2 + \frac{4}{\alpha\beta\Delta} \|V\|^2 + \frac{C^2}{\alpha\beta\Delta} (R/2 - r_1)\ell(R),$$

where

$$(14) \quad \ell(R) := \alpha(\beta - 2\eta) + \frac{2(m+1)\eta - (m+2)\beta}{m} R.$$

Indeed, (9) and (10) give

$$Pr_2 + d = \frac{C}{\alpha\beta} \ell(R).$$

Suppose first that  $R$  can be chosen so that

$$(15) \quad (R/2 - r_1)\ell(R) > 0.$$

For this choice of  $R$ , all three terms in (13) are nonnegative, and the last term is positive unless  $C = 0$ . Therefore

$$C = 0, \quad V = 0, \quad c_y = 0 \quad (y \in X_2).$$

Then (8), (9), and (10) give  $a = 0$ ,  $P = 0$ , and  $d = 0$ . Equation (11) gives  $c_y = 0$  for every  $y \in X$ . The relation (3) then reduces to

$$b\sigma + c\rho = 0.$$

Since  $P = b + c = 0$ , this forces  $b = c = 0$ . The dependence relation is trivial, a contradiction.

It remains to consider the case where no real  $R$  satisfies (15). The polynomial  $(R/2 - r_1)\ell(R)$  is a product of two linear polynomials in  $R$ . It is not identically zero: the first factor is not zero, and  $\ell$  cannot vanish identically, since that would imply both  $\beta = 2\eta$  and  $2(m+1)\eta = (m+2)\beta$ , hence  $\eta = 0$ . Therefore, if it is never positive, the two factors must have the same zero. Hence

$$(16) \quad \ell(2r_1) = 0.$$

Now set  $R = 2r_1$ . Then the last term in (13) is zero, so

$$\sum_{y \in X_2} c_y^2 + \frac{4}{\alpha\beta\Delta} \|V\|^2 = 0.$$

Thus

$$(17) \quad c_y = 0 \quad (y \in X_2), \quad V = 0, \quad a = 0.$$

If  $C = 0$ , then (9) and (10) give  $P = d = 0$ , and (11) again gives  $c_y = 0$  for every  $y \in X$ , leading to a trivial relation as above. Hence  $C \neq 0$ .

Let  $n_1 := |X_1|$ . Since the coefficients on  $X_2$  vanish and  $a = 0$ , equation (11) shows that all coefficients  $c_y$  with  $y \in X_1$  are equal. As their sum is  $C$ ,

$$(18) \quad c_y = \frac{C}{n_1} \quad (y \in X_1).$$

Equation (6) gives

$$(19) \quad \sum_{y \in X_1} y = 0.$$

Moreover, substituting (18) into (11) gives

$$\frac{C}{n_1} + Pr_1 + d = 0.$$

With  $R = 2r_1$ , equations (9) and (10) give

$$Pr_1 + d = \frac{C}{\alpha\beta} B,$$

where

$$B = r_1 \left( 2\eta - \frac{2(m+2)}{m} r_1 \right) + 2\eta r_1 + \alpha\beta - 2(\alpha + \beta)\eta + 2\eta^2.$$

Using (16) to eliminate  $\beta$ , this simplifies to

$$B = -\frac{\beta}{\eta m} D,$$

where

$$D := (m+4)\alpha^2 - 4(m+2)\alpha\eta + 4(m+1)\eta^2.$$

Therefore

$$(20) \quad n_1 = \frac{\alpha\eta m}{D}.$$

Fix  $y \in X_1$ , and let  $k_y$  be the number of points  $y' \in X_1 \setminus \{y\}$  such that

$$|y - y'|^2 = \beta.$$

Since the squared distances inside  $X$  are  $\alpha$  and  $\beta$ , equation (19) gives

$$(n_1 - 1)\alpha + k_y(\beta - \alpha) = \sum_{y' \in X_1, y' \neq y} |y - y'|^2 = 2n_1 r_1.$$

Solving for  $k_y$  gives

$$k_y = \frac{2n_1 r_1 - (n_1 - 1)\alpha}{\beta - \alpha}.$$

The denominator in the next display is nonzero. Indeed, if

$$2(m + 2)\eta = (m + 4)\alpha,$$

then substituting this into  $\ell(2r_1)$  gives

$$-\frac{(m + 4)\alpha^2}{(m + 2)^2},$$

contrary to (16). Hence (16) is equivalent to

$$\beta = \frac{2\eta((m + 2)\alpha - 2(m + 1)\eta)}{(m + 4)\alpha - 2(m + 2)\eta}.$$

Using this identity together with (20), the preceding expression simplifies to

$$(21) \quad k_y = \frac{\alpha(\eta - \alpha)(2(m + 2)\eta - (m + 4)\alpha)^2}{D^2}.$$

The denominator is strictly positive. Indeed, as a quadratic polynomial in  $\eta$ , the discriminant of  $D$  is

$$16\alpha^2((m + 2)^2 - (m + 1)(m + 4)) = -16m\alpha^2 < 0,$$

and its leading coefficient is positive. Also  $0 < \eta < \alpha$ , and the square factor in (21) is nonzero by the preceding paragraph. Therefore (21) gives  $k_y < 0$ , impossible because  $k_y$  is a cardinality. This contradiction excludes  $|X| = G(m)$ , and the lemma follows.  $\square$

## 6. THE TWO-POINT BLOCK CASE

We now complete the analysis of the alternative left by Proposition 4.1. Suppose

$$S = \{p, q\} \cup X \subset \mathbb{R}^d,$$

where  $X$  is a two-distance set contained in the perpendicular bisector  $H$  of the segment  $pq$ , and  $\dim_{\text{aff}} X = d - 1$ . Let  $o$  be the midpoint of  $pq$ , identify  $H$  with  $\mathbb{R}^{d-1}$  using  $o$  as the origin, and write

$$p = (0, \dots, 0, h), \quad q = (0, \dots, 0, -h), \quad h > 0.$$

Put  $n := d - 1$ .

If  $X$  determines at most one nonzero distance, then  $|X| \leq n + 1 = d$ , and therefore

$$|S| \leq d + 2 \leq L(d) + 1 \leq s(d) + 1.$$

Thus we may assume that  $X$  determines exactly two squared distances

$$0 < \alpha < \beta.$$

For  $x \in X$ , put

$$t(x) := |p - x|^2 = |q - x|^2 = |x|^2 + h^2.$$

**Lemma 6.1.** *Among the values  $t(x)$ , there is at most one value outside  $\{\alpha, \beta\}$ .*

*Proof.* Suppose  $x, y \in X$  and  $t(x) \neq t(y)$ . The triangle  $pxy$  is isosceles, and its squared side lengths are

$$t(x), \quad t(y), \quad |x - y|^2.$$

Since  $t(x) \neq t(y)$ , the third squared side length must be one of  $t(x)$  and  $t(y)$ . But  $|x - y|^2 \in \{\alpha, \beta\}$ . Hence two distinct values of  $t$  cannot both lie outside  $\{\alpha, \beta\}$ .  $\square$

**Lemma 6.2.** *If some value  $t(x)$  lies outside  $\{\alpha, \beta\}$ , then*

$$|S| \leq M(d).$$

*Proof.* By Lemma 6.1, there is a unique value  $\gamma \notin \{\alpha, \beta\}$  attained by  $t$  outside  $\{\alpha, \beta\}$ . Put

$$C := \{x \in X : t(x) = \gamma\}.$$

If  $C = X$ , then  $X$  is spherical in  $H$  with center  $o$ , so

$$|S| = |X| + 2 \leq s(n) + 2 \leq s(n) + 3 = s(d - 1) + 3.$$

Assume now that  $C \neq X$ .

Let  $c \in C$  and  $y \in X \setminus C$ . Since  $t(c) = \gamma \notin \{\alpha, \beta\}$  and  $t(y) \in \{\alpha, \beta\}$ , the isosceles triangle  $pcy$  forces

$$(22) \quad |c - y|^2 = t(y).$$

Thus every point of  $X \setminus C$  is the center of a sphere containing  $C$ .

If  $|C| = 1$ , say  $C = \{c\}$ , then (22) gives

$$|c - y|^2 = |y|^2 + h^2 \quad (y \in X \setminus C).$$

Equivalently,

$$c \cdot y = \frac{|c|^2 - h^2}{2} \quad (y \in X \setminus C).$$

Hence  $X \setminus C$  lies in a hyperplane of  $H$ . Therefore

$$|X| \leq 1 + G(n - 1),$$

and so

$$|S| \leq 3 + G(n - 1) \leq L(d) + 1 \leq s(d) + 1.$$

It remains to consider  $|C| \geq 2$ . Put  $b := \dim_{\text{aff}} C$ . Since  $C$  lies on the sphere about  $o$  of squared radius  $\gamma - h^2$ , it is spherical, and

$$|C| \leq s(b) \leq U(b).$$

Moreover, because every point of  $X \setminus C$  is the center of a sphere containing  $C$ , the set  $X \setminus C$  lies in an affine subspace of  $H$  of dimension at most  $n - b$ .

If  $b = n$ , this affine subspace has dimension zero, so  $|X \setminus C| \leq 1$ . Hence

$$|S| \leq s(n) + 3 = s(d - 1) + 3.$$

If  $1 \leq b < n$ , then

$$|X \setminus C| \leq G(n - b),$$

and therefore

$$|X| \leq U(b) + G(n - b) = G(n) - b(n - b) \leq G(n) - 1.$$

It follows that

$$|S| = |X| + 2 \leq G(n) + 1 = L(d) + 1 \leq s(d) + 1.$$

This proves the lemma.  $\square$

**Proposition 6.3.** *In the two-point block case,*

$$|S| \leq M(d).$$

*Proof.* By Lemma 6.2, we may assume

$$(23) \quad t(x) \in \{\alpha, \beta\} \quad (x \in X).$$

Define

$$X_1 := \{x \in X : t(x) = \alpha\}, \quad X_2 := \{x \in X : t(x) = \beta\}.$$

If one of  $X_1, X_2$  is empty, then  $X$  is spherical in  $H$  with center  $o$ , and hence

$$|S| \leq s(n) + 2 \leq s(d-1) + 3.$$

If one of  $X_1, X_2$  has exactly one point, then the other is a spherical two-distance subset of  $H$ . Hence

$$|S| \leq s(n) + 3 = s(d-1) + 3.$$

Finally, suppose that both  $X_1$  and  $X_2$  have at least two points. Applying Lemma 5.1 in dimension  $m = n = d - 1$  gives

$$|X| \leq G(n) - 1.$$

Therefore

$$|S| = |X| + 2 \leq G(n) + 1 = L(d) + 1 \leq s(d) + 1.$$

In every case,  $|S| \leq M(d)$ . □

## 7. PROOF OF THE MAIN THEOREM

*Proof of Theorem 1.1.* Let  $d \geq 2$ , and let  $S \subset \mathbb{R}^d$  be an isosceles set which is not a two-distance set. By Proposition 4.1, either  $|S| \leq M(d)$  or  $S$  lies in the two-point block case. In the latter case, Proposition 6.3 again gives  $|S| \leq M(d)$ . Thus

$$h(d) \leq M(d) = \max\{s(d) + 1, s(d-1) + 3\}.$$

Now let  $S \subset \mathbb{R}^d$  be any isosceles set. If  $S$  is a two-distance set, then  $|S| \leq g(d)$ . If  $S$  is not a two-distance set, then the preceding paragraph gives  $|S| \leq M(d)$ . Hence

$$f(d) \leq \max\{g(d), s(d) + 1, s(d-1) + 3\}.$$

The reverse inequality is Proposition 3.1. Therefore

$$f(d) = \max\{g(d), s(d) + 1, s(d-1) + 3\}.$$

Finally,  $f(1) = 3$ : three equally spaced points on a line are isosceles, while any four distinct points on a line contain three points which are not equally spaced and hence form a non-isosceles triple. This completes the proof. □

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